

From well-formed substring tables to active charts

- Well-formed substring tables: store complete analyses
- But: combination of complete sub-analyses is redone each time

vp → *v-ditr np pp-to*

vp → *v-ditr np np*

- Idea: also store partial results, so that the chart contains
 - passive items: complete results
 - active items: partial results

Representing active chart items

- well-formed substring entry:
 $\text{chart}(i, j, A)$: from i to j there is a constituent of category A
- More elaborate data structure needed to store partial results instead of category:
 - rule considered + how far processing has succeeded
 - dotted rule:
$${}_i[A \rightarrow \alpha \bullet_j \beta] \quad \text{with } A \in N \text{ and } \alpha, \beta \in (\Sigma \cup N)^*$$
- active chart entry: $\text{chart}(i, j, \text{state}(A, \beta))$

Dotted rule examples

- A dotted rule represents a state in processing a rule.
- Each dotted rule is a hypothesis:

	We found a <i>vp</i> if we still find
<hr/> <i>vp</i> → • <i>v-ditr</i> <i>np</i> <i>pp-to</i>	a <i>v-ditr</i> , a <i>np</i> , and a <i>pp-to</i>
<i>vp</i> → <i>v-ditr</i> • <i>np</i> <i>pp-to</i>	a <i>np</i> and a <i>pp-to</i>
<i>vp</i> → <i>v-ditr</i> <i>np</i> • <i>pp-to</i>	a <i>pp-to</i>
<i>vp</i> → <i>v-ditr</i> <i>np</i> <i>pp-to</i> •	nothing

The three actions in Earley's algorithm

In $_i[A \rightarrow \alpha \bullet_j B\beta]$ we call B the *active constituent*.

- **Prediction:** Search all rules realizing the active constituent.
- **Scanning:** Scan over each word in the input string.
- **Completion:** Combine an active edge with each passive edge covering its active constituent.

A closer look at the three actions

Prediction: for each $_i[A \rightarrow \alpha \bullet_j B \beta]$ in chart
for each $B \rightarrow \gamma$ in rules
add $_j[B \rightarrow \bullet_j \gamma]$ to chart

Scanning: let $w_1 \dots w_j \dots w_n$ be the input string
for each $_i[A \rightarrow \alpha \bullet_{j-1} w_j \beta]$ in chart
add $_i[A \rightarrow \alpha w_j \bullet_j \beta]$ to chart

Completion (fundamental rule of chart parsing):

for each $_i[A \rightarrow \alpha \bullet_k B \beta]$ and $_k[B \rightarrow \gamma \bullet_j]$ in chart
add $_i[A \rightarrow \alpha B \bullet_j \beta]$ to chart

Eliminating scanning

Scanning: for each $i[A \rightarrow \alpha \bullet_{j-1} w_j \beta]$ in chart
add $i[A \rightarrow \alpha w_j \bullet_j \beta]$ to chart

Completion: for each $i[A \rightarrow \alpha \bullet_k B \beta]$ and $k[B \rightarrow \gamma \bullet_j]$ in chart
add $i[A \rightarrow \alpha B \bullet_j \beta]$ to chart

Observation: Scanning = completion + words as passive edges

One can thus **replace scanning by:** for each w_j in $w_1 \dots w_n$
add $_{j-1}[w_j \rightarrow \bullet_j]$ to chart

Note: Different from scanning, this is done once for each word, i.e., it is not triggered by every new edge added to the chart.

Earley's algorithm

General setup:

apply prediction and completion to every item added to chart

Start: add ${}_0[start \rightarrow \bullet_0 s]$ to chart

for each w_j in $w_1 \dots w_n$

 add ${}_{j-1}[w_j \rightarrow \bullet_j]$ to chart

Success state: ${}_0[start \rightarrow s \bullet_n]$

A tiny example grammar

- $s \rightarrow np\ vp$
- $np \rightarrow det\ n$
- $vp \rightarrow left$
- $det \rightarrow the$
- $n \rightarrow man$

An example run

- | | |
|---------------------|---|
| start | 1. 0 [start \rightarrow \bullet_0 s] |
| predict from 1 | 2. 0 [s \rightarrow \bullet_0 np vp] |
| predict from 2 | 3. 0 [np \rightarrow \bullet_0 det n] |
| predict from 3 | 4. 0 [det \rightarrow \bullet_0 the] |
| scan "the" | 5. 0 [the \rightarrow \bullet_1] |
| complete 4 with 5 | 6. 0 [det \rightarrow \bullet_1] |
| complete 3 with 6 | 7. 0 [np \rightarrow det \bullet_1 n] |
| predict from 7 | 8. 1 [n \rightarrow \bullet_1 boy] |
| predict from 7 | 9. 1 [n \rightarrow \bullet_1 girl] |
| scan "boy" | 10. 1 [boy \rightarrow \bullet_2] |
| complete 8 with 10 | 11. 1 [n \rightarrow boy \bullet_2] |
| complete 7 with 11 | 12. 0 [np \rightarrow det n \bullet_2] |
| complete 2 with 12 | 13. 0 [s \rightarrow np \bullet_2 vp] |
| predict from 13 | 14. 2 [vp \rightarrow \bullet_2 left] |
| scan "left" | 15. 2 [left \rightarrow \bullet_3] |
| complete 14 with 15 | 16. 2 [vp \rightarrow left \bullet_3] |
| complete 13 with 16 | 17. 0 [s \rightarrow np vp \bullet_3] |
| complete 1 with 17 | 18. 0 [start \rightarrow s \bullet_3] |

The Earley algorithm in PROLOG (earley.pl)

```
% Data structures: chart(From,To,Category)
:- dynamic chart/3.
```

```
% Operator for grammar rules
:- op(1200,xfx,'--->').
```

```
% recognize(+WordList,?Startsymbol)
% top-level predicate for Earley recognizer

recognize(String) :-
    retractall(chart(_,_,_)),
    enter_edge(0,0,state(start,[s])),
    foreach(get_word(Word,JminOne,J,String),
            enter_edge(JminOne,J,state(Word,[]))),
    length(String,N),
    chart(0,N,state(start,[])).
```

```
% enter_edge(+FromIndex,+ToIndex,+Contents)

% a) only add if it does not yet exist:
enter_edge(I,J,State) :- chart(I,J,State), !.

% b) add to chart and try prediction/completion
enter_edge(I,J,State) :-
    assertz(chart(I,J,State)),
    predict(I,J,State),
    complete(I,J,State).
```

```

% prediction(+StartPos,+DotPos,+State)
predict(_,J,state(_, [Cat|_])) :- !,
    foreach((Cat ----> RHS),
            enter_edge(J,J,state(Cat,RHS))).
predict(_,_,_).

```

```

% completion(+StartPos,+DotPos,+State)
complete(K,J,state(B, [])) :- !,
    foreach(chart(I,K,state(A, [B|Beta])),
            enter_edge(I,J,state(A,Beta))).
complete(_,_,_).

```

```
% get_word(-Element,-JminOne,J,+List)
get_word(X,0,1,[X|_]).
get_word(X,JminOne,J,[_|L]) :-
    get_word(X,_,JminOne,L),
    J is JminOne+1.
```

```
% foreach(+Goal1,+Goal2)
foreach(X,Y) :- X, Y, fail.
foreach(_,_).
```

The tiny example grammar (earley/earley_grammar.pl)

```
% lexicon:  
vp ----> [left].  
det ----> [the].  
n ----> [boy].  
n ----> [girl].  
  
% syntactic rules:  
s ----> [np, vp].  
np ----> [det, n].
```

The example run in Prolog

```
| ?- recognize([the,boy,left]).
START:                1: 0-state(start,[s])----0
PRED s in 1:          2: 0-state(s,[np,vp])----0
PRED np in 2:         3: 0-state(np,[det,n])---0
PRED det in 3:        4: 0-state(det,[the])----0
SCAN 1 (the):         5: 0-state(the,[])-----1
COMP 4 + 5:           6: 0-state(det,[])-----1
COMP 3 + 6:           7: 0-state(np,[n])-----1
PRED n in 7:          8: 1-state(n,[boy])-----1
PRED n in 7:          9: 1-state(n,[girl])-----1
SCAN 2 (boy):        10: 1-state(boy,[])-----2
COMP 8 + 10:          11: 1-state(n,[])-----2
COMP 7 + 11:          12: 0-state(np,[])-----2
COMP 2 + 12:          13: 0-state(s,[vp])-----2
PRED vp in 13:       14: 2-state(vp,[left])----2
SCAN 3 (left):       15: 2-state(left,[])-----3
COMP 14 + 15:        16: 2-state(vp,[])-----3
COMP 13 + 16:        17: 0-state(s,[])-----3
COMP 1 + 17:         18: 0-state(start,[])-----3
SUCCESS: 18
```


Improving the efficiency of lexical access

- In the setup just described
 - words are stored as passive items so that
 - prediction is used for preterminal categories. The set of predicted words for a preterminal can be huge.
- If each word in the grammar is introduced by a preterminal rule such as $cat \rightarrow word$ or $lex(cat, word)$, one can
 - add a **passive item for each preterminal category** which can dominate the word instead of for the word itself.

Code change for preterminals as passive edges (earley/passive_preterminals/earley_passive.pl)

In recognize/1 change

```
foreach(get_word(Word, JminOne, J, String),  
        enter_edge(JminOne, J, state(Word, []))),
```

to take into account the preterminal category:

```
foreach((get_word(Word, JminOne, J, String),  
        lex(Cat, Word)),  
        enter_edge(JminOne, J, state(Cat, []))),
```

Towards more flexible control

The algorithms, we saw

- use the Prolog database to store the chart and
- Prolog backtracking on edges in chart instead of an explicit agenda.

Alternatively, one can

- explicitly introduce an **agenda**
- to store and work off edges in any order one likes.

Top-down recognizer with agenda

```
% Data structures: chart(From,To,state(LHS,RHS-LIST))

% Operator for grammar rules
:- op(1200,xfx,'--->').

% recognize(+WordList)
% top-level predicate for Earley recognizer

recognize(String) :-
    enter_string(String,0,N,Agenda),
    FullAgenda = [chart(0,0,state(start,[s])) | Agenda],
    fill_chart(FullAgenda,[],Chart),
    element(chart(0,N,state(start,[],[])),Chart).
```

```
enter_string([],N,N,[]).
enter_string([Word|RestString],JminOne,N,FullAgenda) :-
    J is JminOne + 1,
    findall(chart(JminOne,J,state(Cat,[])),
            lex(Cat,Word),
            FirstAgenda),
    enter_string(RestString,J,N,AgendaRest),
    append(FirstAgenda,AgendaRest,FullAgenda).
```

```
% enter_edges(+EdgeList,+ChartIn,-ChartOut)

enter_edges([],X,X).
enter_edges([Edge | Edges],ChartIn,ChartOut) :-
    enter_edge(Edge,ChartIn,ChartMid),
    enter_edges(Edges,ChartMid,ChartOut).

enter_edge(Edge,Chart,Chart) :- member(Edge,Chart),!.
enter_edge(Edge,Chart,[Edge|Chart]).
```

```

% fill_chart(+Agenda,+ChartIn,-ChartOut)

fill_chart([],X,X).
fill_chart([Edge|RestAgenda],ChartIn,ChartOut) :-
    ChartMid = [Edge|ChartIn],
    %
    predict(Edge,PredictAgenda),
    complete(Edge,ChartMid,CompleteAgenda),
    %
    append(PredictAgenda,RestAgenda,RestPredAgenda),
    append(RestPredAgenda,CompleteAgenda,NewAgenda),
    %
    fill_chart(NewAgenda,ChartMid,ChartOut).

```

```
predict(chart(_, J, state(_, [B|_])), Agenda) :- !,  
    findall(chart(J, J, state(B, Gamma)),  
        (B ----> Gamma),  
        Agenda).  
predict(_, []).
```

```
complete(chart(K, J, state(B, [])), Chart, Agenda) :- !,  
    findall(chart(I, J, state(A, Beta)),  
        element(chart(I, K, state(A, [B|Beta])), Chart),  
        Agenda).  
complete(_, _, []).
```



```
% element(?Element,+List)
```

```
element(X,[X|_]).
```

```
element(X,[_|L]) :-  
    element(X,L).
```

```
% append(+List,?List,-List) or append(-List,?List,+List)
```

```
append([],L,L).
```

```
append([H|T],L,[H|R]) :-  
    append(T,L,R).
```