

Basic DCG notation for encoding CFGs

A DCG rule has the form “*LHS* --> *RHS*.” with

- *LHS*: a Prolog atom encoding a non-terminal, and
- *RHS*: a comma separated sequence of
 - Prolog atoms encoding non-terminals
 - Prolog lists encoding terminals

Examples for some context free grammar rules:

- $S \rightarrow NP VP$
`s --> np, vp.`
- $S \rightarrow NP \text{ thinks } S$
`s --> np, [thinks], s.`
- $S \rightarrow NP \text{ picks up } NP$
`s --> np, [picks, up], np.`
- $S \rightarrow NP \text{ picks } NP \text{ up}$
`s --> np, [picks], np, [up].`
- $NP \rightarrow \epsilon$
`np --> [].`

More complex terms in DCGs

Non-terminals can be any Prolog term, e.g.:

`s --> np(Per, Num) ,
 vp(Per, Num) .`

`s(s_node(NP, VP)) --> np(NP) ,
 vp(VP) .`

Restriction:

- The *LHS* has to be a non-variable, single term (plus possibly a sequence of terminals).

Additional notation for the RHS of DCGs

The *RHS* can include

- **disjunctions** expressed by the “;” operator, e.g.:

```
vp --> vintr;  
      vtrans, np.
```

- **groupings** are expressed using parenthesis “()”
- **extra conditions** in the form of prolog relation calls enclosed in “{ }”, e.g.:

```
s --> np(Case), vp,  
      {check_case(Case)}.
```

```
s --> {write('rule 1'), nl},  
      np,  
      {write('after np'), nl},  
      vp,  
      {write('after vp'), nl}.
```

- the **cut** “!” (can occur without enclosing “{}”).

Meta-variables

On the *RHS*, variables can be used for non-terminals and terminals, i.e. as meta-variables. E.g.:

```
verb([up]) --> [pick].
```

```
vp --> verb(Particle),      % pick
        np,                  % the ball
        Particle.           % up
```

Restriction:

- The value of the variable has to be known at the time Prolog attempts to prove the subgoal represented by the variable.

Towards a basic DCG for English

X-bar Theory

Generalizing over possible phrase structure rules, one can attempt to specify DCG rules fitting the following general pattern:

$$X^2 \rightarrow \text{specifier}^2 X^1$$
$$X^1 \rightarrow X^1 \text{ modifier}^2$$
$$X^1 \rightarrow \text{modifier}^2 X^1$$
$$X^1 \rightarrow X^0 \text{ complement}^{2*}$$

To turn this general X-bar pattern into actual DCG rules,

- X has to be replaced by one of the atoms encoding syntactic categories, and
- the bar-level needs to be encoded as an argument of each predicate encoding a syntactic category.

Noun, preposition, and adjective phrases

Some example rules

```
n(2,Num) --> pronoun(Num) .
n(2,Num) --> proper_noun(Num) .
n(2,Num) --> det(Num), n(1,Num) .
n(2,plur) --> n(1,plur) .
n(1,Num) --> pre_mod, n(1,Num) .
n(1,Num) --> n(1,Num), post_mod .
n(1,Num) --> n(0,Num) .
...

p(2,Pform) --> p(1,Pform) .
p(1,Pform) --> adv, p(1,Pform) .
                % slowly past the window
p(1,Pform) --> p(0,Pform), n(2,_).
...

a(2) --> deg, a(1). % very simple
a(1) --> adv, a(1). % commonly used
a(1) --> a(0).
```

Verb phrases and sentences

Some example rules

$v(2, Vform, Num) \rightarrow v(1, Vform, Num) .$
 $v(1, Vform, Num) \rightarrow adv,$
 $\quad v(1, Vform, Num) .$
 $v(1, Vform, Num) \rightarrow v(1, Vform, Num),$
 $\quad verb_postmods .$
 $v(1, Vform, Num) \rightarrow v(0, intrans, Vform, Num) .$
 $v(1, Vform, Num) \rightarrow v(0, trans, Vform, Num),$
 $\quad n(2) .$
 $v(1, Vform, Num) \rightarrow v(0, ditrans, Vform, Num),$
 $\quad n(2),$
 $\quad n(2) .$

...

 $s(Vform) \rightarrow n(2, Num),$
 $\quad v(2, Vform, Num) .$